

Searching

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Topics in Economic Theory

Overview

1. Pandora's Problem
2. Martingales and etc.
3. Gittins-Jones Index
4. Pricing with Pandora Consumers
5. Summary

Overview

1. Pandora's Problem

- Setup
- Ordering Alternatives
- Optimal Stopping
- Optimal Search
- Variations
- Application: R&D and Project Selection

2. Martingales and etc.

3. Gittins-Jones Index

4. Pricing with Pandora Consumers

5. Summary

Directed Search/Info Acquisition

Before: when to stop.

Now: which info to acquire.

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Setup (Weitzman 1979 Ecta)

$T < \infty$ alternatives, $t \in [T] = \{1, \dots, T\}$ and an outside option 0;

Payoffs $X_t \sim F_t$, independent;

DM can pay cost c_t to learn X_t ;

Recall: DM can stop at any time and pick best alternative explored thus far.

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Motivation:

Firm interviewing applicants;

Consumer searching for a product.

What changes?

Given an order of search, problem is *exactly* the same as before.

Main difficulty: deciding what to learn about next.

Simplifying assumption: independence of payoffs across alternatives

⇒ focus on *history-independent search orders* WLOO (why?).

Ordering Alternatives

Feasible orderings Π :

$$\Pi := \{\pi \in \mathbb{N}_0^{\mathbb{N}_0} \mid \pi \text{ is bijective and } \pi(t) = t \ \forall t \notin [T]\}.$$

π : permutation of elements in $[T] = \{1, \dots, T\}$; defines an order of search.

π makes use only of time t .

$\pi(t)$: alternative DM searches/learns about if they haven't yet stopped by time t .

Ordering Alternatives

Fixing order π , problem as before: i.e. **optimal stopping and choosing**:

$X_{\pi(t)} \equiv X_t^\pi$ as the (gross) payoff associated to alternative $\pi(t) = n$;

$c_{\pi(t)}$ associated search cost;

$M_t^\pi := \max_{s \leq t} X_s^\pi$: highest (gross) payoff thus far;

$Y_t^\pi := M_t^\pi - \sum_{s \leq t} c_{\pi(s)}$, $t \leq T$; $X_t^\pi = Y_t^\pi = -\infty$ for $t > T$; $X_0 = Y_0$ given;

$\mathbb{T}^\pi$: stopping times taking values in $\mathbb{N}_0$ adapted to natural filtration given $\{X_t^\pi\}$.

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Pandora's Problem

$$\max_{\pi \in \Pi} \max_{\tau \in \mathbb{T}^\pi} \mathbb{E}[Y_\tau^\pi] \quad (W)$$

Optimal Stopping

Leveraging what we know:

$\forall \pi$, there is optimal stopping time; Π finite $\implies$ there is solution.

WLOO focus on earliest optimal stopping time for each order.

Solving Pandora's Problem

Start with simple stopping rules:

If problem were monotone, it'd suffice to consider 1-sla/threshold policy

$$\bar{x}_t := \inf\{x \in \mathbb{R} \mid x \geq \mathbb{E}[(X_t \vee x)] - c_t\}.$$

Whether problem monotone or not depends on order of search π .

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Given π , define following stopping time:

$$\tau^\pi := \min \left\{ t \geq 0 \mid M_t^\pi \geq \max_{s>t} \bar{x}_{\pi(s)} \right\}.$$

Solving Pandora's Problem

If DM stops at t it better be that

$$Y_t^\pi \geq \mathbb{E}[Y_{t+1}^\pi \mid \mathcal{F}_t^\pi] \iff M_t^\pi \geq \mathbb{E}[M_t^\pi \vee X_{\pi(t+1)} \mid \mathcal{F}_t^\pi] - c_{\pi(t+1)},$$

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i.e. DM can choose to search not only $X_{\pi(t+1)}$ but any of remaining unsearched alternatives, $X_{\pi(t+h)}$, $h > 0$.

If they're optimally stopping, it must not be profitable to continue and try out any of the remaining alternatives.

Solving Pandora's Problem

Under optimal π and (earliest) optimal stopping time τ , on $\{\tau = t\}$ (stopping at t)

$$M_t^\pi - \sum_{s \leq t} c_{\pi(s)} \geq \mathbb{E}[M_t^\pi \vee X_{t+h}^\pi] - \sum_{s \leq t} c_{\pi(s)} - c_{\pi(t+h)}, \quad \forall h > 0$$

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Compare to $\tau^\pi := \min \left\{ t \geq 0 \mid M_t^\pi \geq \max_{s > t} \bar{x}_{\pi(s)} \right\}$.

Implies $\tau \geq \tau^\pi$.

Natural conjecture to check: τ^π is optimal.

Pandora's Optimal Stopping Time

Proposition

Suppose (π, τ) solve Pandora's problem. Then, $\mathbb{E}[Y_{\tau}^{\pi}] = \mathbb{E}[Y_{\tau^{\pi}}^{\pi}]$, τ^{π} is regular, and $\tau^{\pi} \leq \tau$.

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Contradicts regularity of τ . Hence, $\{t \geq \tau^\pi\} \cap \{\tau > t\} = \emptyset$ and $\{\tau^\pi \leq t\} = \{\tau \leq t\}$.

Conclude: $\tau^\pi = \tau$. More: $\forall$ optimal stopping time τ' , $\mathbb{E}[Y_{\tau'}^\pi] = \mathbb{E}[Y_\tau^\pi] = \mathbb{E}[Y_{\tau^\pi}^\pi]$ and $\tau' \geq \tau^\pi$.

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Regularity of τ^π follows from definition

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If for all t , $\bar{x}_{\pi(t)} \geq \bar{x}_{\pi(t+1)}$ problem would be monotone, since

$$M_t^\pi \geq \max_{s>t} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+1)} \implies Y_t^\pi \geq \mathbb{E}[Y_{t+1}^\pi \mid \mathcal{F}_t^\pi]$$

Then τ^π would simply correspond to the 1-sla rule!

Pandora's Optimal Search Order

Claim: There's beauty and order in the universe. Should check for simple solutions.

Conjecture: if π were *not* inducing a decreasing sequence of $\bar{x}_{\pi(t)}$, rearranging π so that $\{\bar{x}_{\pi(t)}\}_t$ is decreasing sequence improves payoffs.

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If (π, τ^π) solve Pandora's problem, then $\{\bar{x}_{\pi(s)}\}_{s \leq t}$ is nonincreasing for all $t : \mathbb{P}(\tau^\pi \geq t) > 0$

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Iterating argument optimal order π sat. $\{\bar{x}_{\pi(s)}\}_{s \leq t}$ nonincreasing.

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We find out that the solution to Pandora's problem is s.t. (1) optimal order is simply to order alternatives according to $\bar{x}_{\pi(t)}$; (2) use 1-sla stopping rule!

Variations to Pandora's Problem

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Choosing without Search

Things go a bit awry, but Doval (2018 JET) provides sufficient conditions for Pandora's rule to remain optimal except for last alternative.

Also shows how for simple setting (binary outcomes) breaking independence leads to spectacular failure of Pandora's rule to be optimal.

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E.g. alternatives as different components of research project.

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Flexible Learning

I haven't checked thoroughly, but I think an open problem is to retain independence and characterise the optimal solution when allowing for more flexible learning.

Hope for something nice/tractable that would do well in applications.

E.g., when assessing candidates, scan some CVs, then dig into some good ones and interview them; if turn out not great, go back to pile.

Application: R&D and Project Selection

T different research projects

Project n :

- Ex-ante prob success $p_n \in (0, 1)$ and payoff $r_n > 0$; ow payoff 0.
- Associated cost c_n .
- DM only considers projects with positive exp. value: $p_n r_n - c_n > 0$.
- Outside option yields 0.

Application: R&D and Project Selection

T different research projects

Project n :

- Ex-ante prob success $p_n \in (0, 1)$ and payoff $r_n > 0$; ow payoff 0.
- Associated cost c_n .
- DM only considers projects with positive exp. value: $p_n r_n - c_n > 0$.
- Outside option yields 0.

Solution

$$\bar{x}_n := \inf\{x \in \mathbb{R} \mid x \geq \mathbb{E}[r_n \vee x] - c_n\} = r_n - \frac{c_n}{p_n},$$

DM trades off expected net reward and prob success:

- Explore first projects with high potential reward net of cost scaled by prob success.
- Stop myopically (1-sla).

Overview

1. Pandora's Problem

2. Martingales and etc.

- Martingales
- Doob's Decomposition
- Doob's Martingale Convergence and Optional Stopping
- Wald's Equations

3. Gittins-Jones Index

4. Pricing with Pandora Consumers

5. Summary

Martingales and etc.

We'll do a quick detour to recall some fundamental results about martingales.

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Martingales and etc.

Let $X = \{X_t\}_{t \in [T]}$ be an adapted process.

X is a **supermartingale** if $\mathbb{E}[X_t^-] < \infty \forall t \in T$ and $\mathbb{E}[X_t | \mathcal{F}_s] \leq X_s$ a.s. for all $s \leq t$, $s, t \in [T]$.

X is a **submartingale** if $-X$ is a supermartingale.

X is a **martingale** if it is both a super- and submartingale. Specifically, a martingale satisfies $\mathbb{E}[X_t | \mathcal{F}_s] = X_s$ for $s \leq t$.

X is a **predictable process** if X_t is $\mathcal{F}_{t-1}$ -measurable.

Doob's Decomposition Theorem

Doob's Decomposition Theorem

If X is an $\mathbb{F}$ -adapted process satisfying $\mathbb{E}[\sup_{s \geq t} |X_s|] < \infty$, then $\exists$ martingale $Z : Z_0 = 0$ and an integrable predictable process $A : A_0 = 0$ s.t. $X_t = X_0 + Z_t + A_t \quad \forall t$, with decomposition being unique a.s.

Implication: super/submartingale = martingale + a.s. de/increasing predictable process.

Doob's Martingale Convergence Theorem

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If X is submartingale s.t. $\sup_{t \geq 0} \mathbb{E}[X_t^+] < \infty$, then X_t converges pointwise to X_∞ , $X_t(\omega) \rightarrow X_\infty(\omega)$, and $X_\infty < \infty$ a.s.

Moreover, if X is martingale, it is uniformly integrable if and only if X_t converges a.s. and in L^1 to X_∞ satisfying $X_t = \mathbb{E}[X_\infty \mid \mathcal{F}_t]$ for all t .

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Recall, (i) X is uniformly integrable if $\lim_{a \rightarrow \infty} \sup_t \mathbb{E}[|X_t| \mathbf{1}_{\{|X_t| > a\}}] = 0$; (ii)

$$X_t \xrightarrow{L^1} X_\infty \iff \mathbb{E}[|X_t - X_\infty|] \rightarrow 0.$$

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Doob's Optional Stopping Theorem

Let X be supermartingale (resp. submartingale) and τ a stopping time. Suppose one of the following holds:

- (i) $\exists c < \infty : \tau \leq c$ a.s.
- (ii) $\mathbb{E}[\tau] < \infty$ and $\exists c < \infty : \forall t, \mathbb{E}[|X_{t+1} - X_t| \mid \mathcal{F}_t] \leq c$ on $\{\tau > t\}$.
- (iii) $\exists c < \infty : \forall t, |X_{t \wedge \tau}| \leq c$ a.s.

Then, X_τ is a.s. well-defined r.v. and $\mathbb{E}[X_\tau] \leq$ (resp. $\geq$) $\mathbb{E}[X_0]$.

Doob's Martingale Convergence Theorem

Counterexample

Wealth after sequence of iid fair bets, $S_t = \sum_{s \leq t} X_s$, where $X_s = \pm 1$ wp $1/2$ and $X_0 = 0$.
 $\tau := \inf\{t \geq 0 \mid S_t = 1\}$. Then S_τ is martingale.

$\mathbb{E}[\tau] = \infty$ and optional stopping theorem doesn't apply: $\mathbb{E}[S_\tau] = 1 > 0 = \mathbb{E}[S_t]$.

S_t doesn't converge in mean. Also, $\{S_t\}$ not u.i.

Wald's Equation

Let X be s.t. (i) $\mathbb{E}[X_t] < \infty$, (ii) $\forall t \mathbb{E}[X_t \mathbf{1}\{\tau \geq t\}] = \mathbb{E}[X_t] \mathbb{P}(\tau \geq t)$, and (iii) $\sum_{t \in \mathbb{N}} \mathbb{E}[|X_t| \mathbf{1}\{\tau \geq t\}] < \infty$. Define $S_\tau := \sum_{t=1}^\tau X_t$ and $T_\tau := \sum_{t=1}^\tau \mathbb{E}[X_t]$.

Then $\mathbb{E}[S_\tau] = \mathbb{E}[T_\tau] < \infty$. If, moreover, $\mathbb{E}[X_t] = m \forall t$ and $\mathbb{E}[\tau] < \infty$, then $\mathbb{E}[S_\tau] = \mathbb{E}[\tau]m$.

Wald's Equations

There are a number of versions of this result typically labeled Wald's First/ Second/ Third Identity which follow from the optional stopping theorem:

Wald's First/Second/Third Identities

Let $\{X_t\}_{t \in \mathbb{N}}$ be a stochastic process such that X_t are independent, with $\mathbb{F}$ denoting its natural filtration, and τ be an adapted stopping time with $\mathbb{E}[\tau] < \infty$. Define (i) $S_t := \sum_{\ell=1}^t X_\ell$; (ii) $m_t := \sum_{\ell=1}^t \mathbb{E}[X_\ell]$; (iii) $v_t := \sum_{\ell=1}^t \mathbb{V}(X_\ell)$; (iv) $\phi(\theta) := \mathbb{E}[\exp(\theta X_1)]$; (v) $M_t^1 := S_t - m_t$; (vi) $M_t^2 := (S_t - m_t)^2 - v_t$; and (vii) $M_t^3 := \phi(\theta)^{-t} \exp(\theta S_t)$.

1. If $\sup_t \mathbb{E}[|X_t|] < \infty$, then M_t^1 is a martingale and $\mathbb{E}[M_\tau^1] = \mathbb{E}[M_1^1] = 0$. In particular, if X_t are iid with mean $\mathbb{E}[X_t] = m$, then $\mathbb{E}[S_\tau] = m\mathbb{E}[\tau]$.
2. If $\sup_t \mathbb{E}[X_t^2] < \infty$, then M_t^2 is a martingale and $\mathbb{E}[M_\tau^2] = \mathbb{E}[M_1^2] = 0$. In particular, if X_t are iid with variance $\mathbb{V}(X_t) = \sigma^2$, then $\mathbb{E}[(S_\tau - m\tau)^2] = \sigma^2\mathbb{E}[\tau]$.
3. If X_t are iid, the moment generating function $\phi(\theta) < \infty$, and τ is a.s. bounded or $M_t^3 \mathbf{1}\{\tau \geq t\} \leq \delta < \infty$ for all t , then M_t^3 is a martingale and $\mathbb{E}[\phi(\theta)^{-\tau} \exp(\theta S_\tau)] = 1$.

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Multi-Armed Bandit Problems

Setup:

Each period $t = 0, 1, \dots$,

DM ('gambler') chooses one action ('arm') $a_t \in A := \{1, \dots, K\}$

receives a random payoff $x_t^{a_t}$

whose distribution $F^{a_t}(\cdot; s_t^k)$ depends on the state $s_t = (s_t^1, \dots, s_t^K)$,

and the state evolves according to a Markov chain s.t. $s_{t+1}^k = \phi_k(x_t^k, s_t^k)$ if $a_t = k$ and $s_{t+1}^k = s_t^k$ if $a_t \neq k$. ('restless' bandit when allowing $s_{t+1}^k \neq s_t^k$ even if $a_t \neq k$)

Payoffs are discounted by $\delta \in [0, 1)$.

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Motivation: learning-by-doing + choosing what to do.

Choosing and switching streaming platform subscription.

Learning about job match value and decide to switch to another job.

Learning about effects of enacted policy.

Joining a queue and switching to another another.

Draw-down retirement savings by depleting different assets.

Seller learning demand by experimenting with prices.

Scheduling of experiments (e.g., Pandora's boxes!)

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Example of Learning Framing:

Unknown parameters: $\theta^k \in \Theta^k$, $\theta^k \sim \mu_0^k \in \Delta(\Theta^k)$.

$\theta = (\theta^1, \dots, \theta^K)$ and $\mu_0 = \times_k \mu_0^k$ (independence; product measure).

Objective payoff distributions: $X_t^k \sim G^k(\cdot \mid \theta^k)$ iid.

Posterior beliefs: $\mu_{t+1}^k = \mu_t^k \mid X_t^k$ if $a_t = k$ and $\mu_{t+1}^k = \mu_t^k$ if $a_t \neq k$ (our Markovian 'state').

Subjective payoff distributions: $F^k(\cdot \mid \mu_t^k) := \mathbb{E}_{\mu_t^k}[G^k(\cdot \mid \theta^k)]$.

Actions entail payoffs *and* learning about payoff distribution.

Multi-Armed Bandit Problems

Setup:

Histories: $h^t := (s_0, s_1, \dots, s_t) \in H_t$ and $\mathcal{H} := \cup_{t=0}^{\infty} H_t$.

Strategies: $\alpha \in A^{\mathcal{H}}$.

Payoffs $\mathbb{E}[\sum_{t=0}^{\infty} \delta^t X_t^{\alpha_t}]$.

Goal: $V(s_0) := \sup_{\alpha} \mathbb{E}[\sum_{t=0}^{\infty} \delta^t X_t^{\alpha_t}]$.

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Denote $X^k(s_t^k) \sim F^k(\cdot; s_t^k)$.

Theorem (Gittins & Jones, 1974)

The optimal policy satisfies $\alpha(s_t) \in \arg \max_{a \in A} m^a(s_t^a)$, where

$$(1 - \delta)m^k(s_0^k) := \sup_{\tau} \frac{\mathbb{E}[\sum_{t=0}^{\tau-1} \delta^k X^k(s_t^k)]}{\mathbb{E}[\sum_{t=0}^{\tau-1} \delta^k]}.$$

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Comments:

Decomposes K -dimensional problem to solve K 1-dimensional problems.

Formulation very general: with transition can capture use-costs of arm, countable number of arms, etc.

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Seminal papers: Gittins & Jones (1974, 1979), Gittins (1979), Weber (1992).

Also Gittins (1989), Karatzas (1984; Brownian bandits), Banks & Sundaram (1994; switching costs).

Applications: Pricing and learning demand: Rothschild (1974); McLennan (1984); Rustichini & Wolinsky (1995); Keller & Rady (1999); Bergemann & Välimäki (2006), Bonatti (2010). Strategic experimentation: Bolton & Harris (1999); Keller, Rady, & Cripps (2005); Strulovici (2010), etc.

Multi-Armed Bandit Problems

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Consider two actions: pull arm k or get lump-sum reward M .

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$$V(s_0^k, M) := \sup_{\tau} \mathbb{E} \left[\sum_{t=0}^{\tau-1} \delta^t X^k(s_t^k) + \delta^{\tau} M \right] \quad (\text{SP})$$

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Note: (i) V increasing in M ,

(ii) convex in M (maximising over linear functions of M), and

(iii) $\exists M' : V(s_0^k, M) = M \iff M \geq M'$ (both stopping and continuation region are intervals).

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Experimentation suboptimal if $M = m^k(s^k)$. Hence, $\forall \tau$,

$$m^k(s_0^k) \geq \mathbb{E} \left[\sum_{t=0}^{\tau-1} \delta^t X^k(s_t^k) + \delta^{\tau} m^k(s_0^k) \right] \iff m^k(s_0^k)(1 - \mathbb{E}[\delta^{\tau}]) \geq \mathbb{E} \left[\sum_{t=0}^{\tau-1} \delta^t X^k(s_t^k) \right].$$

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Since $\sum_{s=0}^{\tau-1} \delta^s = \frac{1-\delta^{\tau}}{1-\delta}$, $m^k(s_0^k)(1 - \mathbb{E}[\delta^{\tau}]) = (1 - \delta)m^k(s_0^k)\mathbb{E}[\sum_{t=0}^{\tau-1} \delta^t]$. Then,

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with equality at optimal stopping:

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Multi-Armed Bandit Problems

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Intuition from Weber (1992):

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- Once arm is abandoned, restart process of lowering the price offer.

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- Operator runs arm until Gittins-Jones index falls below price, i.e. its original Gittins-Jones Index.
- Once arm is abandoned, restart process of lowering the price offer.
- Since operators get zero surplus and they are operating under optimal rules, this method of allocating arms results in maximal surplus to owner, solving the original MAB problem.

Theorem (Gittins & Jones, 1974)

The optimal policy satisfies $\alpha(s_t) \in \arg \max_{a \in A} m^a(s_t^a)$, where

$$(1 - \delta)m^k(s_0^k) := \sup_{\tau} \frac{\mathbb{E}[\sum_{t=0}^{\tau-1} \delta^k X^k(s_t^k)]}{\mathbb{E}[\sum_{t=0}^{\tau-1} \delta^k]}.$$

Very nice result. Hard to compute index in closed-form other than in special cases.

Fragility: result breaks with minor variations.

- non-geometric discounting, e.g., fixed time horizon;
- arms with correlated priors;
- actions affecting more than one arm at a time;
- payoffs depending on state of 2 or more arms;
- delayed feedback;
- 'restless' arms that change state without being pulled;
- switching costs; etc.

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 - Market Equilibrium
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Pricing with Pandora Consumers

Pricing with search (Choi, Dai, & Kim, 2018 Ecta)

Sellers

N sellers; each supplies product n at price $p_n \geq 0$; marginal cost c_n .

p : vector of prices; p_{-n} : prices of n 's competitors.

Demand for n given prices: $D_n(p)$.

Sellers maximise profit: $\pi_n(p) := D_n(p)(p_n - c_n)$.

Pricing with Pandora Consumers

Consumers

Consumer i 's valuation of n : $X_n = V_n + W_n - p_n$.

V_n is known; W_n is idiosyncratic-value component revealed only when consumer learns about n .

Cost to acquire info and learn W_n : $k_n > 0$.

Outside option X_0 ; search with recall.

V_n and W_n are independently drawn from F_n and G_n , both smooth.

Independence allows use of Pandora's rule.

If consumer stops after searching sellers $A \subseteq [N]$, the consumer accrues a payoff

$$\max_{n \in A \cup \{0\}} X_{i,n} - \sum_{m \in A} k_m, \text{ where } p_0 = k_0 = 0.$$

Consumers know p .

Pandora's Rule

Define $\bar{w}_n$: $k_n = \mathbb{E}[(W_n - \bar{w}_n)^+]$

Given V_n define $\bar{x}_n := V_n + \bar{w}_n - p_n$

Given order $\pi \in \Pi$, let $M_t^\pi := \max_{s \leq t} X_{\pi(t)} = \max_{s \leq t} V_{\pi(t)} + W_{\pi(t)} - p_{\pi(t)}$

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Proposition

Given a price vector p and realization V , it is optimal for the consumer

- (1) to learn about sellers in decreasing order of $\bar{x}_n$, with the optimal order of search being given by $\pi \in \Pi$ such that $\bar{x}_{\pi(t)} \geq \bar{x}_{\pi(t+1)}$; and
- (2) to stop whenever $M_t^\pi \geq \bar{x}_{\pi(t+1)}$, with the earliest optimal stopping time being given by $\tau := \min\{t \geq 0 \mid M_t^\pi \geq \bar{x}_{\pi(t+1)}\}$

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Proof

This is just Pandora's rule.

Theorem 1

Given a price vector p and realizations V, W , consumer i chooses product n if $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$ and only if $X_n \wedge \bar{x}_n \geq \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$

We've already proved a version of this for the satisficing setup.

(A proof for this specific setup is [here](#).)

Write consumers' expected payoff given V as $\mathbb{E}[\max_n (X_n \wedge \bar{x}_n) \vee X_0 \mid V]$

Intuition

Let $\bar{M} := \max_{n \in [T]} X_n \wedge \bar{x}_n \vee X_0$:

(1) n is chosen if $X_n \wedge \bar{x}_n > \bar{M}$, and

(2) the consumer learns about n whenever $\bar{x}_n > \bar{M}$, incurring in cost

$$k_n = \mathbb{E}[(X_n - \bar{x}_n)^+].$$

Market Equilibrium

Let (1) H_n be s.t. $X_n \wedge \bar{x}_n \sim H_n$;

(2) $\bar{H}_n$ be s.t. $\max_{m \neq n} X_m \wedge \bar{x}_m \vee X_0 \sim \bar{H}_n$.

Market Equilibrium

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Demand for n :

$$D_n(p) = \mathbb{P}(\max_{m \neq n} X_m \wedge \bar{x}_m \vee X_0 \leq x < X_n \wedge \bar{x}_n) = \int (1 - H_n(x_n)) d\bar{H}_n(x_n).$$

p is **eqm price** if $\forall n$,

$$p_n \in \arg \max_{p'_n} D_n(p'_n, p_{-n})(p'_n - c_n).$$

FOC:

$$\frac{1}{p_n - c_n} = - \frac{dD_n(p)/dp_n}{D_n(p)}.$$

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Assumption 1: H_n and $1 - H_n$ are log-concave $\forall n$.

Assumption 2: $\text{supp } H_n$ has no upper bound $\forall n$.

Theorem

Under Assumption 1, $D_n(p)$ is log-concave in p_n and $\log D_n(p)$ has strictly increasing differences in (p_n, p_m) .

Under Assumptions 1 and 2, there is a unique Nash equilibrium of the pricing game, and this equilibrium is in pure strategies.

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Existence of PSNE from game being supermodular (assumed away).

Uniqueness is due to particular structure of quasilinear preferences.

Beautiful result; ties-in:

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Issue: H_n is endogenous!

Theorem 2

Under Assumption 1, $D_n(p)$ is log-concave in p_n and $\log D_n(p)$ has strictly increasing differences in (p_n, p_m) .

Under Assumptions 1 and 2, there is a unique Nash equilibrium of the pricing game, and this equilibrium is in pure strategies.

Sufficient conditions for uniqueness

- (i) f_n and g_n log-concave and $\sup \text{supp } F_n = +\infty \implies 1 - H_n$ log-concave.
- (ii) F_n entails sufficiently high variance and $(f_n(\underline{v}_n) = 0 \text{ or } \inf \text{supp } F_n = -\infty) \implies H_n$ log-concave.
- (iii) $f'_n(X_0 + c_n - \bar{w}_n) \leq 0 \implies H_n$ log-concave on relevant part of its support.

Horizontal Differentiation

Imperfect Price Competition

Always exist *some* consumers who value n much more than others.

Typically this is *assumed*, e.g. loyal buyers vs shoppers, horizontal differentiation.

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Distribution H_2 is **more dispersed** than H_1 **above** w if $H_2^{-1}(b) - H_2^{-1}(a) \geq H_1^{-1}(b) - H_1^{-1}(a)$ for any $H_1(w) < a \leq b < 1$.

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Higher values more dispersed.

Simplifying assumption: **symmetric environment**; i.e. for all $n, m \in [T]$

$F \equiv F_n = F_m$, $G \equiv G_n = G_m$, and $k \equiv k_n = k_m$ (implying $H \equiv H_n = H_m$), and

$C \equiv C_n = C_m$.

Proposition 3

In symmetric environments, eqm price increases as H becomes more dispersed above $X_0 + c$ and $H(X_0 + c)$ weakly decreases.

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Proposition 4

In symmetric environments, either condition is sufficient for H to become more dispersed above $X_0 + c$:

- (1) G (match values) becomes more dispersed and f is log-concave, or
- (2) F (prior/common values) becomes more dispersed, g is log-concave, f is decreasing above $X_0 + c - \bar{w}$, and $\inf \text{supp } F \leq X_0 + c - \bar{w}$.

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'Effective' (or post-learning) preference heterogeneity.

+ dispersion in $G \implies$ + dispersion 'ex-post' valuations; *and*

+ dispersion in $G \implies$ + value to learning.

Combined, both reinforce + effective heterogeneity.

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In symmetric environments, either condition is sufficient for H to become more dispersed above $X_0 + c$: (...)

(2) F (prior/common values) becomes more dispersed, g is log-concave, f is decreasing above $X_0 + c - \bar{w}$, and $\inf \text{supp } F \leq X_0 + c - \bar{w}$.

Intuition for (2):

F + dispersed $\equiv$ + ex-ante heterogeneity.

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Implies higher ex-post heterogeneity, *but* may also preclude learning.

Learning is source of effective heterogeneity, no conflicting forces.

Conditions in (2) ensure stronger incentives to learning, so that we have both forces going in same direction.

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In symmetric environments, if f is log-concave, both eqm price and profit decrease in cost to learning k .

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Why do we get this then?

Intuition relies on presumption that prices not known to consumers prior to learning.

Here: prices known; it's valuations that are not known.

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When learning costs are higher, consumers less willing to learn and demand + sensitive to prices.

Thus: $\uparrow$ learning costs $\implies$ + price competition $\implies$ $\downarrow$ eqm price $\implies$ $\downarrow$ profits.
(this last bit is not immediate: outside option)

Overview

1. Pandora's Problem
2. Martingales and etc.
3. Gittins-Jones Index
4. Pricing with Pandora Consumers
5. Summary

Summary

Search as flexible framework for directed info acquisition
a lot to be done...

- Directed job search from first principles;

- Market entry with learning;

- Structural estimation of consumer demand;

- Manipulation/obfuscation via search costs;

- Uncertainty about prices and valuations;

- Dynamic pricing;

- Inference based on decision times (see Choi & Smith for comparative statics);

- Optimal R&D funding design.

Searching

Duarte Gonçalves

University College London

Topics in Economic Theory

Proof

(Back)

Suppose $\bar{x}_{\pi(s)}$ nonincreasing for $s \geq t+1$, but $\bar{x}_{\pi(t)} < \bar{x}_{\pi(t+1)}$.

Assume that $\tau^\pi \geq t$ occurs with positive prob., as ow it is WL to rearrange the order following t (no impact on payoffs).

Define new order δ : same as π except it swaps t -th and $(t+1)$ -th alternatives.

Define stopping time $\tau := \min\{s \geq 0 \mid M_s^\delta \geq \max_{n>s} \bar{x}_{\pi(s)}\}$

Then $\{\tau \leq s\} = \{\tau^\pi \leq s\}$, $\forall s \neq t$; i.e. only difference between τ^π and τ is if they disagree in stopping at t .

$\{\tau \in \{t, t+1\}\} = \{\tau^\pi \in \{t, t+1\}\} = \{M_{t-1}^\pi < \bar{x}_{\pi(t+1)} \text{ and } M_{t+1}^\pi \geq \bar{x}_{\pi(t+2)}\}$.

$\{\tau^\pi = t\} = \{M_{t-1}^\pi < \bar{x}_{\pi(t+1)} = \max_{s>t-1} \bar{x}_{\pi(s)} \text{ and } X_t = M_t \geq \bar{x}_{\pi(t+1)} = \max_{s>t-1} \bar{x}_{\pi(s)}\}$,
 whereas $\{\tau = t\} = \{M_{t-1}^\pi < \bar{x}_{\pi(t+1)} = \max_{s>t-1} \bar{x}_{\pi(s)} \text{ and } X_{t+1} = M_t^\delta \geq \bar{x}_{\pi(t+1)} = \max_{s>t-1} \bar{x}_{\pi(s)}\}$.

$$\mathbb{E}[Y_{\tau^\pi}^\pi - Y_{\tau}^\delta] = \mathbb{E}[\mathbf{1}_{\{\tau \in \{t, t+1\}\}}(Y_{\tau^\pi}^\pi - Y_{\tau}^\delta)]$$

Proof

(Back)

$\{\tau \in \{t, t+1\}\} = \{\tau^\pi \in \{t, t+1\}\} = \{M_{t-1}^\pi < \bar{x}_{\pi(t+1)}\} \cap \{M_{t+1}^\pi \geq \bar{x}_{\pi(t+2)}\}$. Furthermore,
 $\{\tau^\pi = t\} = \{M_{t-1}^\pi < \max_{s>t-1} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+1)}\} \cap \{X_t^\pi = M_t^\pi \geq \max_{s>t-1} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+1)}\}$;
 $\{\tau = t\} = \{M_{t-1}^\delta = M_{t-1}^\pi < \max_{s>t-1} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+1)}\} \cap \{X_{t+1}^\pi = M_t^\delta \geq \max_{s>t-1} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+1)}\}$.
 $\{\tau^\pi = t+1\} = \{X_t^\pi \leq M_t^\pi < \max_{s>t} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+1)}\} \cap \{M_{t+1}^\pi = M_{t+1}^\delta \geq \max_{s>t+1} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+2)}\}$;
 $\{\tau = t+1\} = \{X_{t+1}^\pi \leq M_t^\delta < \max_{s>t} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+1)}\} \cap \{M_{t+1}^\pi = M_{t+1}^\delta \geq \max_{s>t} \bar{x}_{\pi(s)} = \bar{x}_{\pi(t+2)}\}$.

Note that

$$\begin{aligned}
 & \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_t^\pi \wedge X_{t+1}^\pi \geq \bar{X}_{\pi(t+1)}\}} (Y_{\tau^\pi}^\pi - Y_\tau^\delta) \\
 &= \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_t^\pi \wedge X_{t+1}^\pi \geq \bar{X}_{\pi(t+1)}\}} ((X_t^\pi - X_{t+1}^\pi) - (c_{\pi(t)} - c_{\pi(t+1)})) \\
 & \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_t^\pi \geq \bar{X}_{\pi(t+1)} > X_{t+1}^\pi\}} (Y_{\tau^\pi}^\pi - Y_\tau^\delta) \\
 &= \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_t^\pi \geq \bar{X}_{\pi(t+1)} > X_{t+1}^\pi\}} ((X_t^\pi - X_{t+1}^\pi) - (c_{\pi(t)} - c_{\pi(t+1)} - c_{\pi(t)})) \\
 &= \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_t^\pi \geq \bar{X}_{\pi(t+1)} > X_{t+1}^\pi\}} c_{\pi(t+1)} \\
 & \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{X}_{\pi(t+1)} > X_t^\pi\}} (Y_{\tau^\pi}^\pi - Y_\tau^\delta) \\
 &= \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{X}_{\pi(t+1)} > X_t^\pi\}} ((X_{t+1}^\pi - X_t^\pi) - (c_{\pi(t)} + c_{\pi(t+1)} - c_{\pi(t+1)})) \\
 &= -\mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{X}_{\pi(t+1)} > X_t^\pi\}} c_{\pi(t)} \\
 & \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{\bar{X}_{\pi(t+1)} > X_t^\pi \vee X_{t+1}^\pi\}} (Y_{\tau^\pi}^\pi - Y_\tau^\delta) \\
 &= \mathbf{1}_{\{\tau \in \{t, t+1\}\}} \mathbf{1}_{\{\bar{X}_{\pi(t+1)} > X_t^\pi \vee X_{t+1}^\pi\}} ((M_{t-1}^\pi - M_{t-1}^\pi) - (c_{\pi(t)} + c_{\pi(t+1)} - c_{\pi(t+1)} - c_{\pi(t)})) = 0
 \end{aligned}$$

Proof

(Back)

Since $1 = 1_{\{X_t^\pi \wedge X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} + 1_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)} > X_{t+1}^\pi\}} + 1_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)} > X_t^\pi\}} + 1_{\{X_t^\pi \vee X_{t+1}^\pi < \bar{x}_{\pi(t+1)}\}}$ we get:

$$\begin{aligned} & \mathbb{E}[Y_{\rho^\pi}^\pi - Y_\tau^\delta] \\ &= \mathbb{E} \left[1_{\{\tau \in \{t, t+1\}\}} \left(1_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} 1_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} ((X_t^\pi - X_{t+1}^\pi) - (C_{\pi(t)} - C_{\pi(t+1)})) \right. \right. \\ & \quad \left. \left. + 1_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)} > X_{t+1}^\pi\}} C_{\pi(t+1)} - 1_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)} > X_t^\pi\}} C_{\pi(t)} \right) \right] \\ &= \mathbb{E} \left[1_{\{\tau \in \{t, t+1\}\}} \left(1_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} 1_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} (X_t^\pi - X_{t+1}^\pi) + 1_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} C_{\pi(t+1)} - 1_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} C_{\pi(t)} \right) \right] \end{aligned}$$

Proof

(Back)

Recalling that $c_{\pi(t)} := \mathbb{E}[(X_t^\pi - \bar{x}_{\pi(t)})^+]$ and $\bar{x}_{\pi(t)} < \bar{x}_{\pi(t+1)}$, we get

$$\begin{aligned}
 & \mathbb{E}[Y_{\rho^\pi}^\pi - Y_\tau^\delta] \\
 &= \mathbb{E} \left[\mathbf{1}_{\{\tau \in \{t, t+1\}\}} \left(\mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} (X_t^\pi - X_{t+1}^\pi) \right. \right. \\
 & \quad \left. \left. + \mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbb{E}[(X_{t+1}^\pi - \bar{x}_{\pi(t+1)})^+] - \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbb{E}[(X_t^\pi - \bar{x}_{\pi(t)})^+] \right) \right] \\
 &< \mathbb{E} \left[\mathbf{1}_{\{\tau \in \{t, t+1\}\}} \left(\mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} (X_t^\pi - X_{t+1}^\pi) \right. \right. \\
 & \quad \left. \left. + \mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbb{E}[(X_{t+1}^\pi - \bar{x}_{\pi(t+1)})^+] - \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbb{E}[(X_t^\pi - \bar{x}_{\pi(t+1)})^+] \right) \right]
 \end{aligned}$$

Proof

(Back)

$$\begin{aligned}
 \mathbb{E}[Y_{\rho^\pi}^\pi - Y_\tau^\delta] &< \mathbb{E} \left[\mathbf{1}_{\{\tau \in \{t, t+1\}\}} \left(\mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} (X_t^\pi - X_{t+1}^\pi) \right. \right. \\
 &\quad \left. \left. + \mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbb{E}[(X_{t+1}^\pi - \bar{x}_{\pi(t+1)})^+] - \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbb{E}[(X_t^\pi - \bar{x}_{\pi(t+1)})^+] \right) \right] \\
 &= \mathbb{E} \left[\mathbf{1}_{\{\tau \in \{t, t+1\}\}} \left(\mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} (X_t^\pi - X_{t+1}^\pi) \right. \right. \\
 &\quad \left. \left. + \mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} (X_{t+1}^\pi - \bar{x}_{\pi(t+1)}) - \mathbf{1}_{\{X_{t+1}^\pi \geq \bar{x}_{\pi(t+1)}\}} \mathbf{1}_{\{X_t^\pi \geq \bar{x}_{\pi(t+1)}\}} (X_t^\pi - \bar{x}_{\pi(t+1)}) \right) \right] \\
 &= 0
 \end{aligned}$$

Theorem

Given a price vector p and realizations V, W , consumer i chooses product n if $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$ and only if $X_n \wedge \bar{x}_n \geq \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.

Theorem

Given a price vector p and realizations V, W , consumer i chooses product n if $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$ and only if $X_n \wedge \bar{x}_n \geq \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.

Proof

[\(Back\)](#)

Only if:

- Case 1: $X_n \wedge \bar{x}_n < \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.
- Case 1a: $X_n \wedge \bar{x}_n = X_n$.

If $X_n < X_0$, consumer will never purchase n .

If $X_n < X_m \wedge \bar{x}_m$ and consumer stops after $\pi^{-1}(m)$, they will never choose n .

If $X_n < X_m \wedge \bar{x}_m$ and stops at $t < \pi^{-1}(m)$, then $M_t^\pi > \bar{x}_m \geq X_m \wedge \bar{x}_m > X_n$ and so they will not choose n .

Theorem

Given a price vector p and realizations V, W , consumer i chooses product n if $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$ and only if $X_n \wedge \bar{x}_n \geq \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.

Proof

(Back)

Only if:

- Case 1: $X_n \wedge \bar{x}_n < \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.
- Case 1b: $X_n \wedge \bar{x}_n = \bar{x}_n < X_n$.

If $\pi^{-1}(m) < \pi^{-1}(n)$, then $M_{\pi^{-1}(n)-1}^\pi \geq (X_m \wedge \bar{x}_m) \vee X_0 > \bar{x}_n$ and so τ calls for stopping and consumer never learns about product n and thus never buys it.

If $\pi^{-1}(m) > \pi^{-1}(n)$ and $(X_m \wedge \bar{x}_m) \geq X_0$, then $X_n > \bar{x}_n \geq \bar{x}_m \geq X_m \wedge \bar{x}_m > X_n \wedge \bar{x}_n = \bar{x}_n$, a contradiction.

If $X_n \wedge \bar{x}_n = \max_m (X_m \wedge \bar{x}_m) < X_0$, then if stop at t (but not earlier) $X_0 \leq M_{s-1}^\pi < \bar{x}_{\pi(s)}$.

This implies $X_0 \geq X_{\pi(s)}$, for all $s = 1, \dots, t$, therefore $M_t^\pi = X_0$.

By assumption, $\bar{x}_n < X_n$, and so if $\pi^{-1}(n) \leq t$, $X_0 > X_n$ and n cannot be chosen.

Theorem

Given a price vector p and realizations V, W , consumer i chooses product n if $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$ and only if $X_n \wedge \bar{x}_n \geq \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.

Proof

[\(Back\)](#)

If:

- Case 2: $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.
- Case 2a: $X_n \geq \bar{x}_n$.

Then $M_{\pi^{-1}(n)} \geq X_n \geq \bar{x}_n \geq \bar{x}_{\pi^{-1}(n)+1}$ and therefore $\tau \leq \pi^{-1}(n)$.

As $\bar{x}_n > X_m \wedge \bar{x}_m$, for all $m : \pi^{-1}(m) < \pi^{-1}(n)$,

$\implies \bar{x}_m \geq \bar{x}_n$ and $X_n \geq \bar{x}_n > X_m$

$\implies n$ must be chosen.

Theorem

Given a price vector p and realizations V, W , consumer i chooses product n if $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$ and only if $X_n \wedge \bar{x}_n \geq \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.

Proof

(Back)

If:

- Case 2: $X_n \wedge \bar{x}_n > \max_{m \neq n} (X_m \wedge \bar{x}_m) \vee X_0$.
- Case 2b: $X_n < \bar{x}_n$.

Then $\bar{x}_m \geq \bar{x}_n > X_n > X_m \wedge x_m$ for all $m : \pi^{-1}(m) < \pi^{-1}(n)$ and $\tau \geq \pi^{-1}(n)$.

If the consumer stops at $t > \pi^{-1}(n)$, then

$\bar{x}_s > X_n > X_s \wedge \bar{x}_s$ for all $s \leq t$ such that $s \neq \pi^{-1}(n)$ and $M_t^\pi = X_n$

$\implies n$ is chosen.